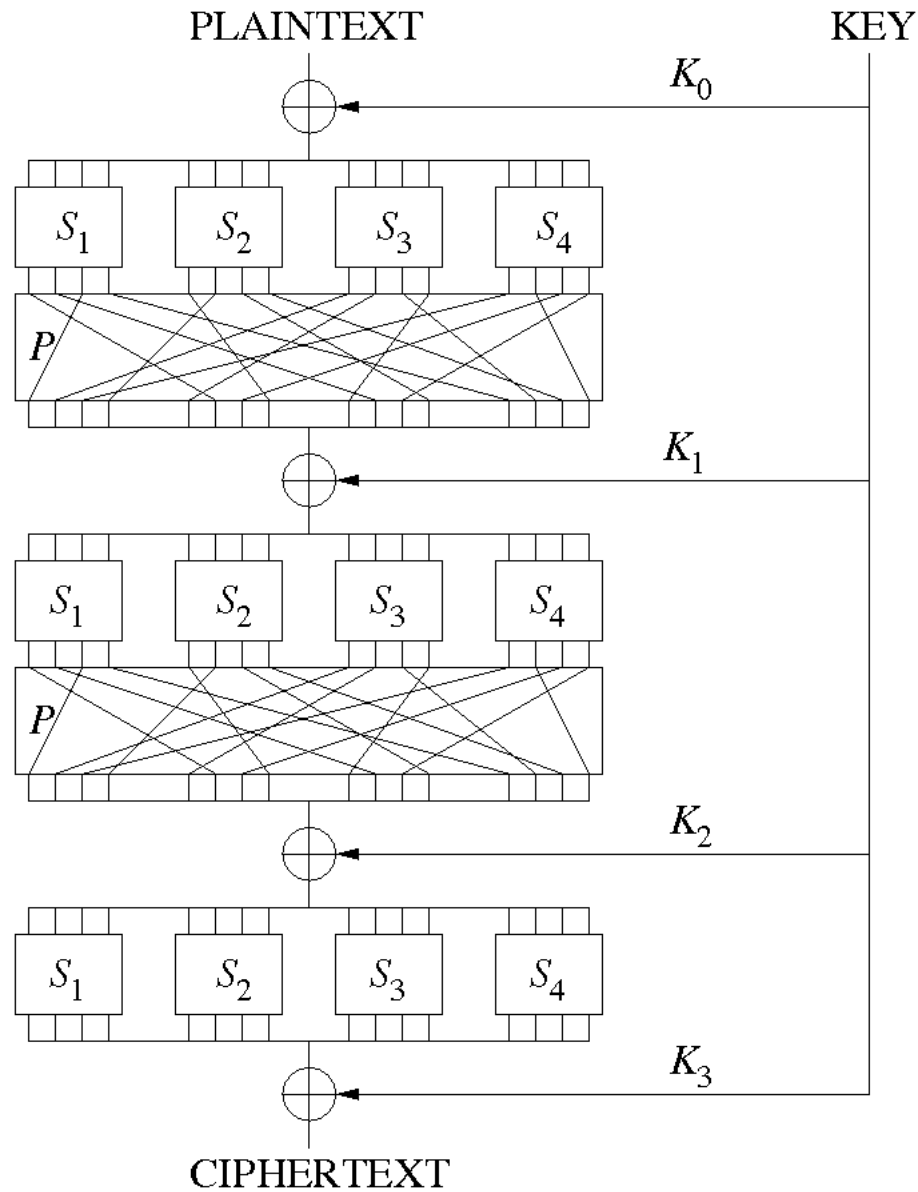


# CBC padding oracle attacks

# CBC padding oracle attack examples

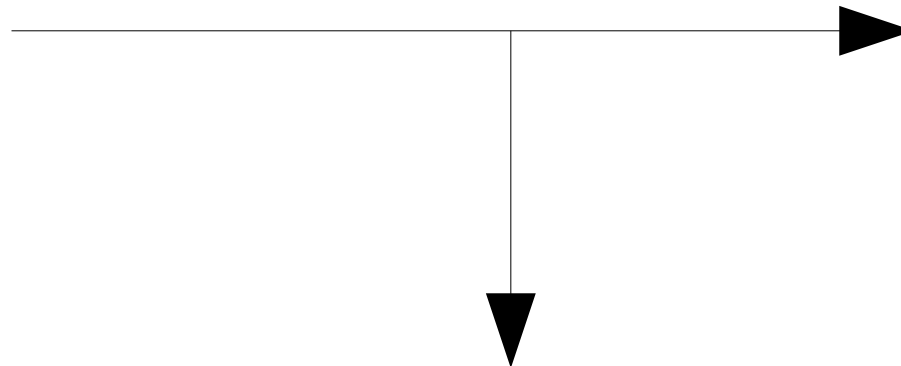
- Serge Vaudenay published the original attack in 2002
  - Applied to web frameworks like Ruby on Rails, ASP.NET, and JavaServer Faces
  - <https://www.iacr.org/cryptodb/archive/2002/EUROCR YPT/2850/2850.pdf>
- POODLE (published by Google in 2014) exploited SSLv3 that is still widely used by web servers and browsers
  - <https://security.googleblog.com/2014/10/this-poodle-bites-exploiting-ssl-30.html>

# Review: AES is a block cipher and is symmetric



Source: Wikipedia

# Alice and Bob have a shared secret key



Eve makes a copy of the ciphertext as it is transmitted from Alice to Bob.

# Alice and Bob have a shared secret key



Eve re-plays modified copies of the encrypted message and learns information about the plaintext from Bob's behavior (e.g., Bob throws an exception for padding error)

# PKCS#7 padding

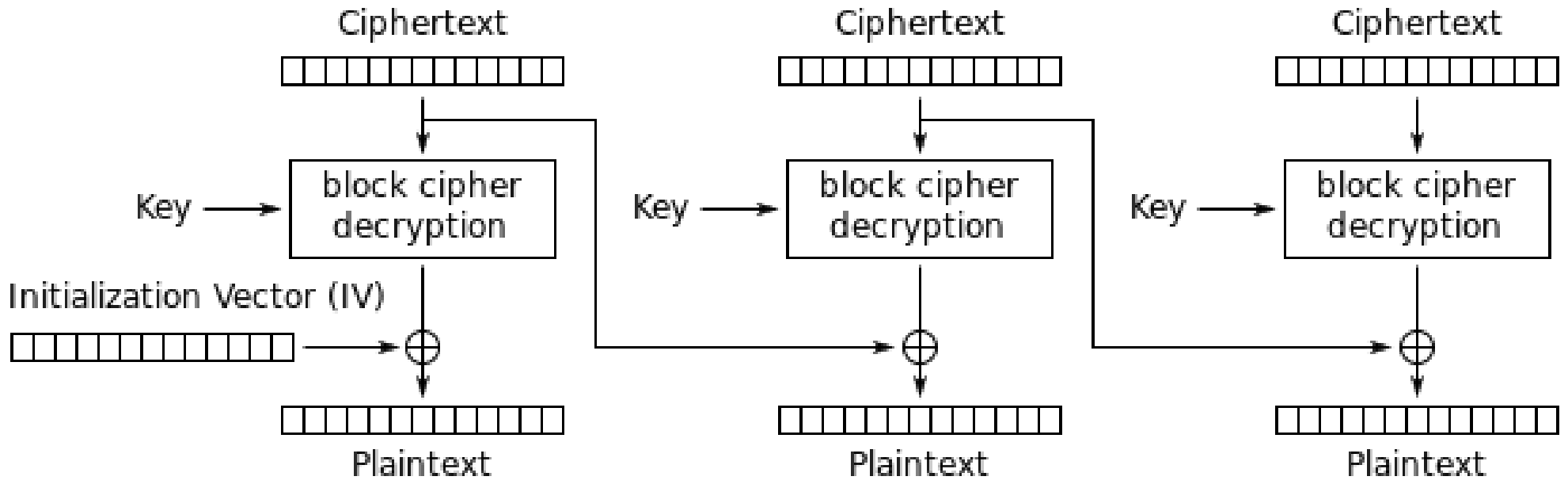
- AES always encrypts in 128-bit blocks
  - 128 bits == 16 bytes
- If you fill up blocks, that's great
  - The last block might not be full
- Need an “unambiguous” way to pad the last block so the decrypting party knows the padding to throw out
  - *E.g.*, PKCS#7 (PKCS == Public Key Cryptography Standards)



# When last block is decrypted

- Check last byte of the last block, that's the number of bytes of padding
  - Call it N
- There should be N N's on the end
  - If not, throw a padding error
  - If so, remove them, they're padding
    - Might remove the whole last block if  $N = 16$  (or 10 in hex)





Cipher Block Chaining (CBC) mode decryption

# Requirements for attack

- Ability to modify ciphertexts and replay them
  - Chosen ciphertext attack
- A padding oracle
  - *i.e.*, something that tells you whether the corresponding plaintext (for any ciphertext you send) has valid padding or not

# Example plaintext (we don't know the plaintext yet before the attack)

H	e	l	l	o	20	W	o	r	l	d	!	\n	03	03	03
---	---	---	---	---	----	---	---	---	---	---	---	----	----	----	----


# Example protocol for a client to send an encrypted message to a server

N	u	m	b	l	k	s	:	1	K	e	y	l	D	:	A3
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
98	CC	BE	01	FF	26	39	97	85	A1	02	1E	BC	A5	7E	98

# Example protocol for a client to send an encrypted message to a server

Number of blocks

Which key?



N	u	m	b	l	k	s	:	1	K	e	y	l	D	:	A3
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
98	CC	BE	01	FF	26	39	97	85	A1	02	1E	BC	A5	7E	98

# Example protocol for a client to send an encrypted message to a server

N	u	m	B	l	k	s	:	1	K	e	y	l	D	:	A3
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
98	CC	BE	01	FF	26	39	97	85	A1	02	1E	BC	A5	7E	98

IV is randomly chosen but visible on the wire and known to you, won't be 0 like in this illustration

# Example protocol for a client to send an encrypted message to a server

N	u	m	B	l	k	s	:	1	K	e	y	l	D	:	A3
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
98	CC	BE	01	FF	26	39	97	85	A1	02	1E	BC	A5	7E	98

Ciphertext is what you want to decrypt, you will recover the plaintext without needing to know the key!

# Server response is visible to you

- “Message decrypted successfully”  
---or---
- “Padding error during decryption”

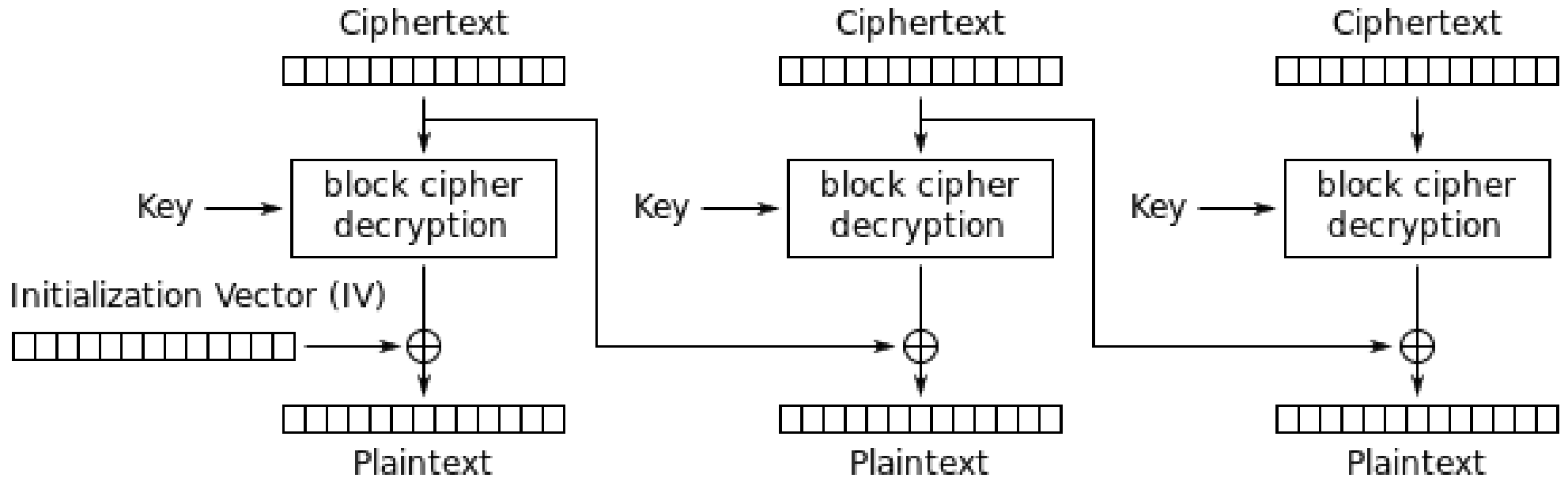


# You can record a client message and replay it to the server

N	u	m	b	l	k	s	:	1	K	e	y	l	D	:	A3
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	00
98	CC	BE	01	FF	26	39	97	85	A1	02	1E	BC	A5	7E	98

Try every value of this byte from 00 to FF





## Cipher Block Chaining (CBC) mode decryption

# Suppose two values give valid padding

- 00 gives valid padding, this is just confirmation that the original plaintext has valid padding
  - 02 also gives valid padding
    - Can recover one byte of plaintext:  
 $Q \text{ XOR } 02 == 01$ , so...  $Q == 01 \text{ XOR } 02 == 03$
- Q is the byte of plaintext we're trying to guess

# WTF?

N	u	m	b	l	k	s	:	1	K	e	y	l	D	:	A3
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	02
98	CC	BE	01	FF	26	39	97	85	A1	02	1E	BC	A5	7E	98



H	e	l	l	o	20	W	o	r	l	d	!	\n	03	03	01
---	---	---	---	---	----	---	---	---	---	---	---	----	----	----	----

“Information only has meaning in that it is subject to interpretation”

$$01 \text{ XOR } 02 = 03$$

N	u	m	b	l	k	s	:	1	K	e	y	l	D	:	A3
00	00	00	00	00	00	00	00	00	00	00	00	00	00	00	01
98	CC	BE	01	FF	26	39	97	85	A1	02	1E	BC	A5	7E	98



H	e	l	l	o	20	W	o	r	l	d	!	\n	03	03	02
---	---	---	---	---	----	---	---	---	---	---	---	----	----	----	----



Now attack here

# Discussion

- You still don't know the key, and probably never will
- It doesn't matter how secure AES is or the size of the key
- CBC is probably the most commonly used mode
- What if a byte is already what it needs to be?
- What if there is more than one block?

# References

- <https://grymoire.wordpress.com/2014/12/05/cbc-padding-oracle-attacks-simplified-key-concepts-and-pitfalls/>